



Blue Gene/P Architecture

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IBM, USA



IBM's Blue Gene® Supercomputer

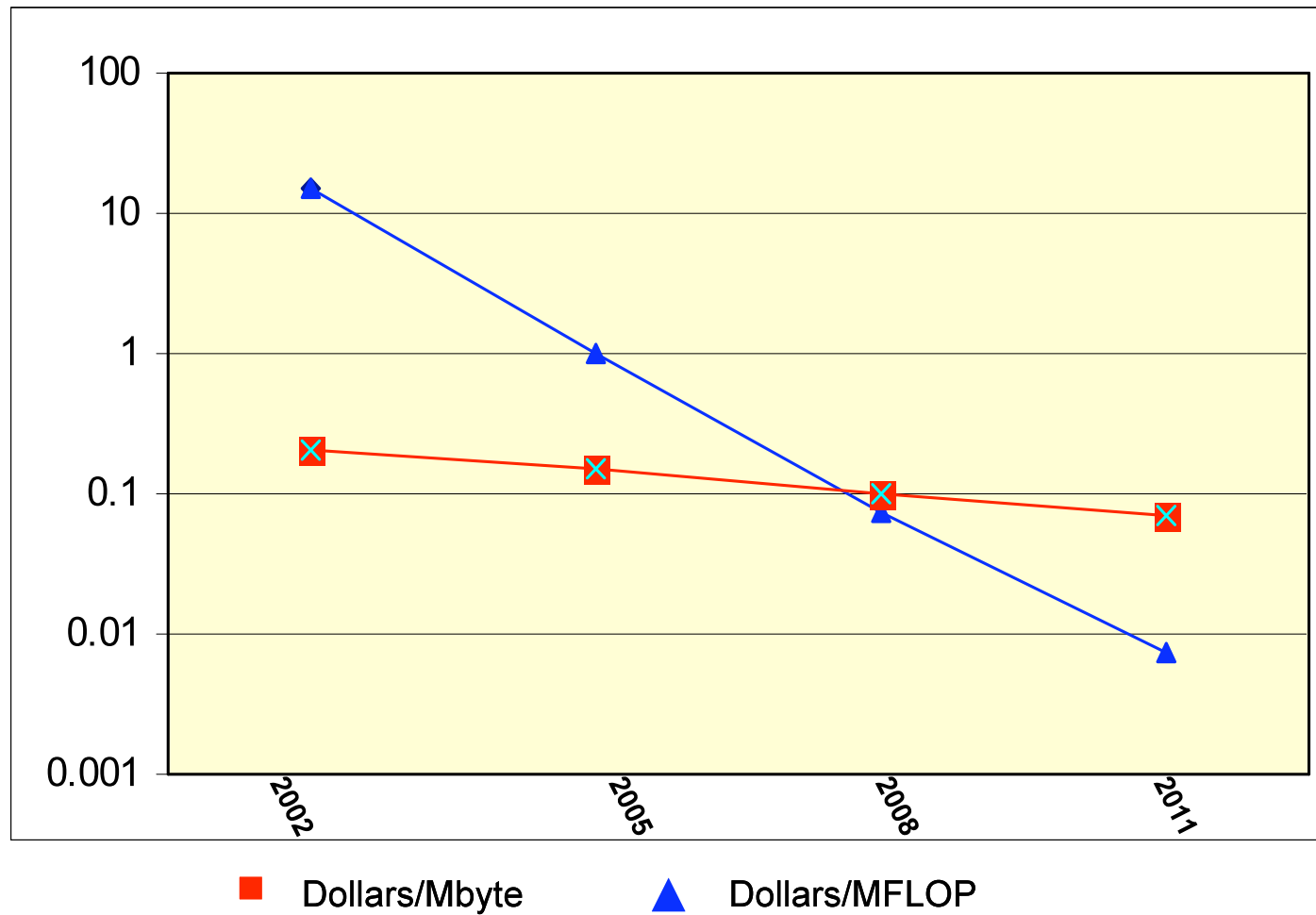
Applying *Innovation that Matters* to High Performance Computing

- Leadership performance, ultra scalability
- Space saving design
- Power efficient computing
- High reliability
- Ease of system management
- Familiar programming methods

Blue Gene/P Architectural Highlights

- Scaled performance relative to BG/L through density and frequency
 - 1.2x from frequency bump 700 MHz => 850 MHz
 - 2x performance through doubling the processors/node
- Enhanced function
 - 4 way SMP, cache coherent, supports threads, OpenMP
 - Improved memory subsystem supporting higher bandwidths
 - DMA for torus, remote put-get, user programmable memory prefetch
 - Memory chip kill implemented.
 - Enhanced performance counters (including 450 core)
 - Architectures of double Hummer FPU, torus, collective network, barrier, and JTAG networks were left intact.
- Higher signaling rate
 - 2.4x higher bandwidth, lower latency for Torus and Tree networks
 - 10x higher bandwidth for Ethernet IO

GFLOPs vs DRAM Price Reductions



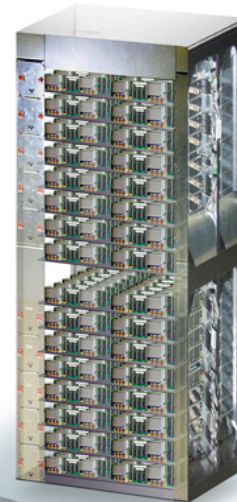
Blue Gene/P

System



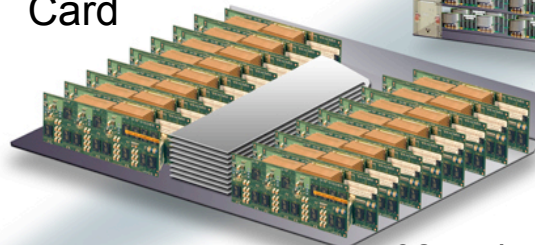
72 Racks
73728 nodes
294912 cores
1 PF
144 TB DDR2

Rack



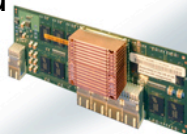
32 node-cards
1024 nodes
4096 cores
13.9 TF
2 TB DDR2

Node
Card



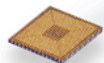
32 nodes
128 cores
435 GF
64GB DDR2

Compute
Card



1 node
4 cores
13.6 GF
2GB DDR2

Chip



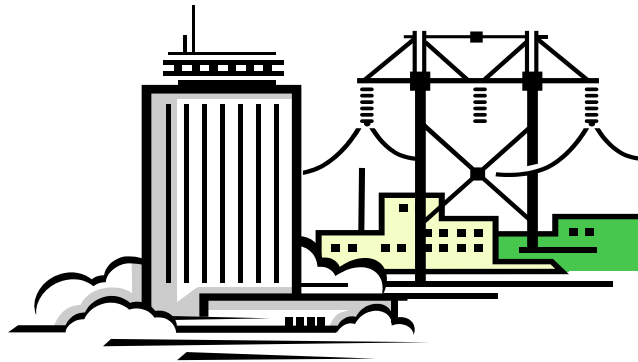
4 cores
13.6 GF

Blue Gene/P Packaging Characteristics

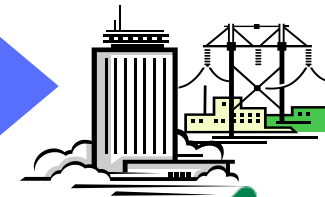
Frequency	850MHz PowerPC 450
First level Package	FC-PBGA, 29 mm × 29 mm; 528s + 255p/g
Compute card	6s-8p FR4 card, 145 mm x 54.6 mm
Node card	6s-10p FR4 card, 427 mm x 400 mm
Midplane	5s-9p FR4 card, 640 mm x 550 mm
Rack	4ft x 3ft x 42U, 5500CFM, 40KW delivered

Application	kW within rack	bulk power efficiency	Total kW Rack Power (480VAC Wall Power)
LINPACK	28.7	0.91	31.5
Ave application	21.3	0.91	23.4
Rack Idle	8.60	0.87	9.9

Blue Gene/L's power efficiency is first rank



Blue Gene requires
75-80% less power and
space than COTS cluster



ENERGYGUIDE

Compare the Energy Use of this Computer
with Others Before You Buy.



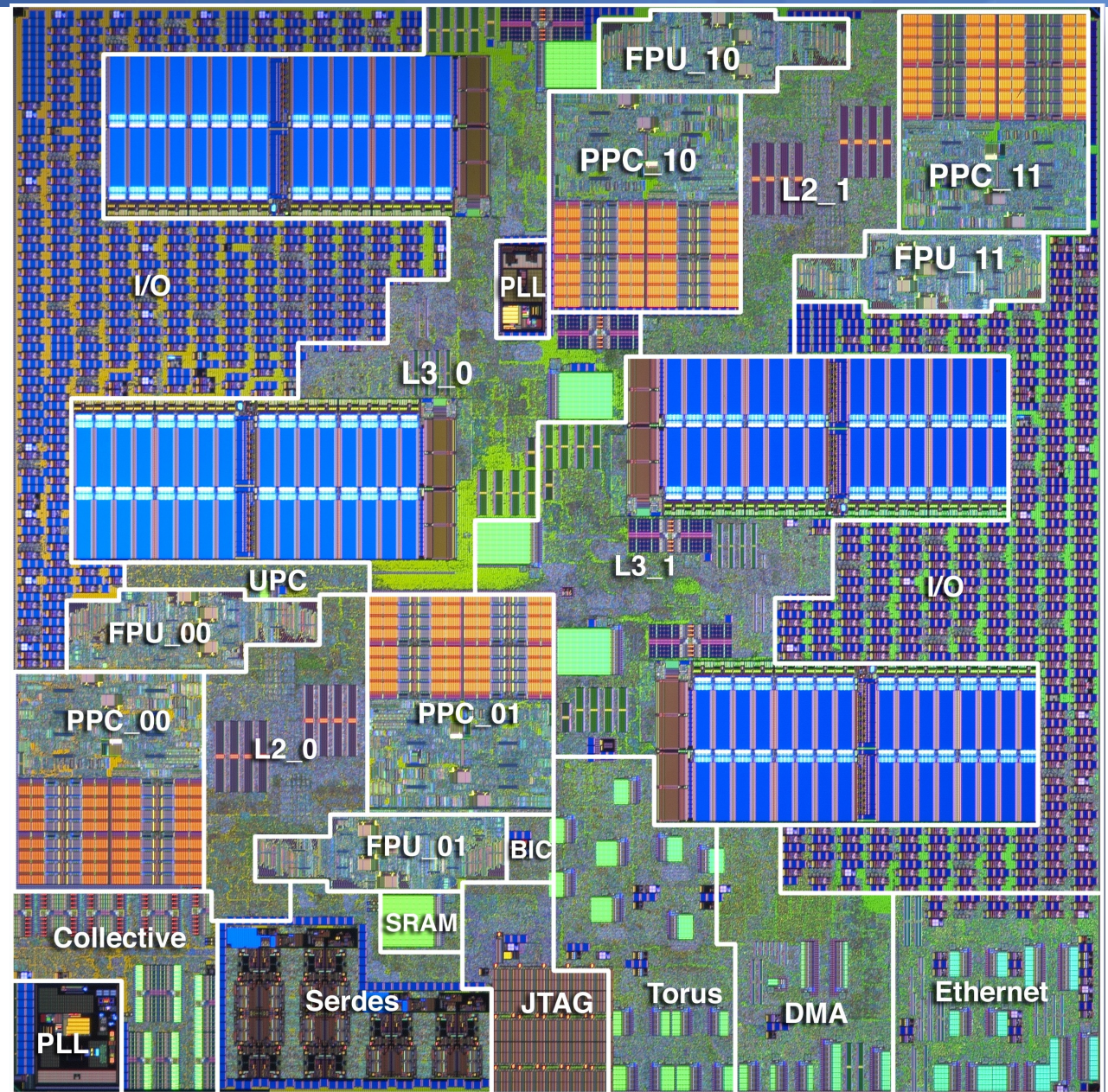
Green500 Rank	Machine	Location	MFlops / KW
1	Cell (IBM)	IBM Germany	488.14
1	Cell (IBM)	Fraunhofer ITWM	488.14
3	Cell (IBM)	DOE/NNSA/LANL	437.43
4	Blue Gene/P (IBM)	ANL	371.75
7	Blue Gene/P (IBM)	ORNL (Eugene)	371.67

Computational efficiency is moving from sustained-to-peak (aka MPH/HP) to performance-per-watt (aka MPG)

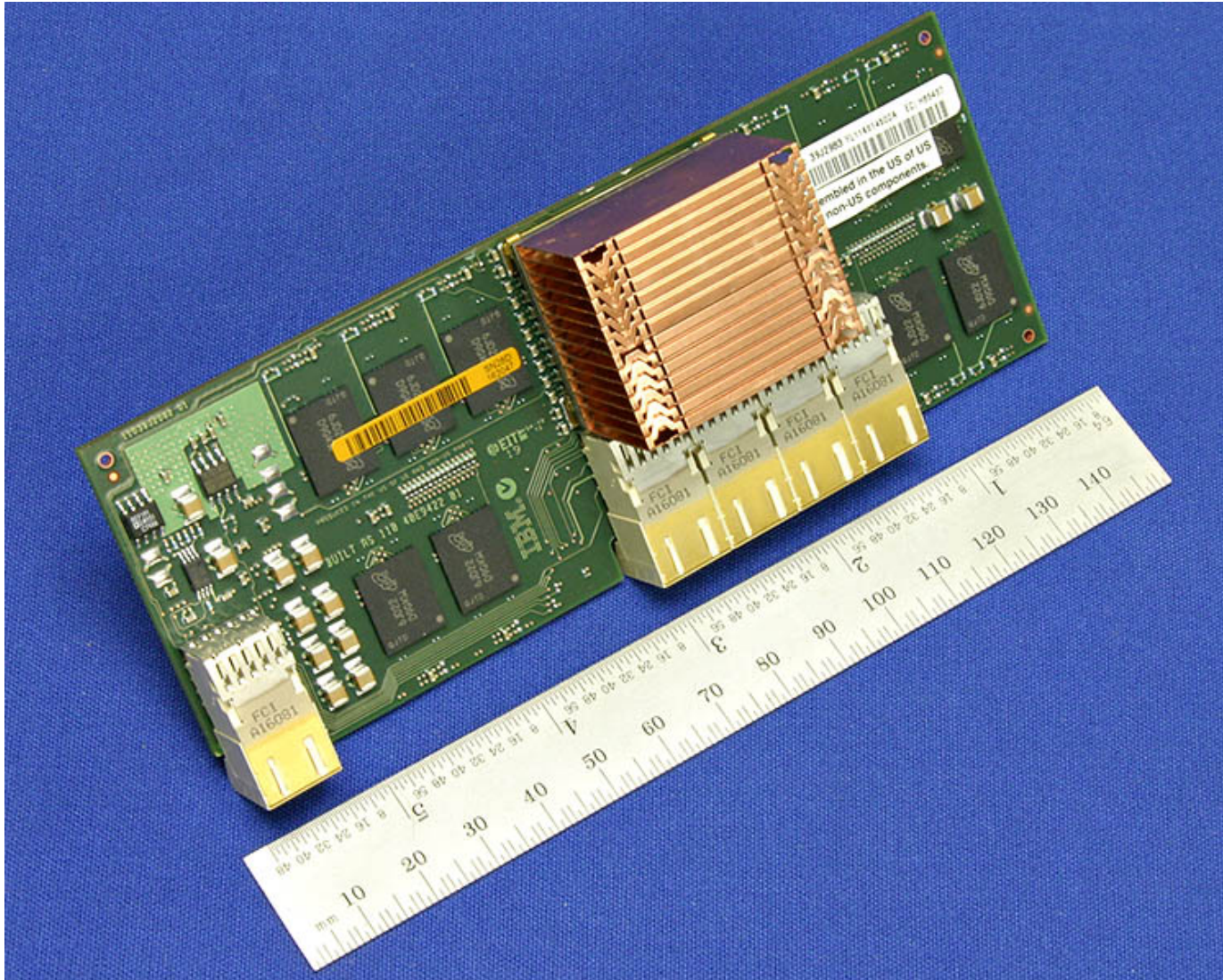
*John Shalf, NERSC/LBNL,
"The Landscape of Computer Architecture" presented at
ISC07, Dresden, June 2007*

BPC chip DD2.1 die photograph

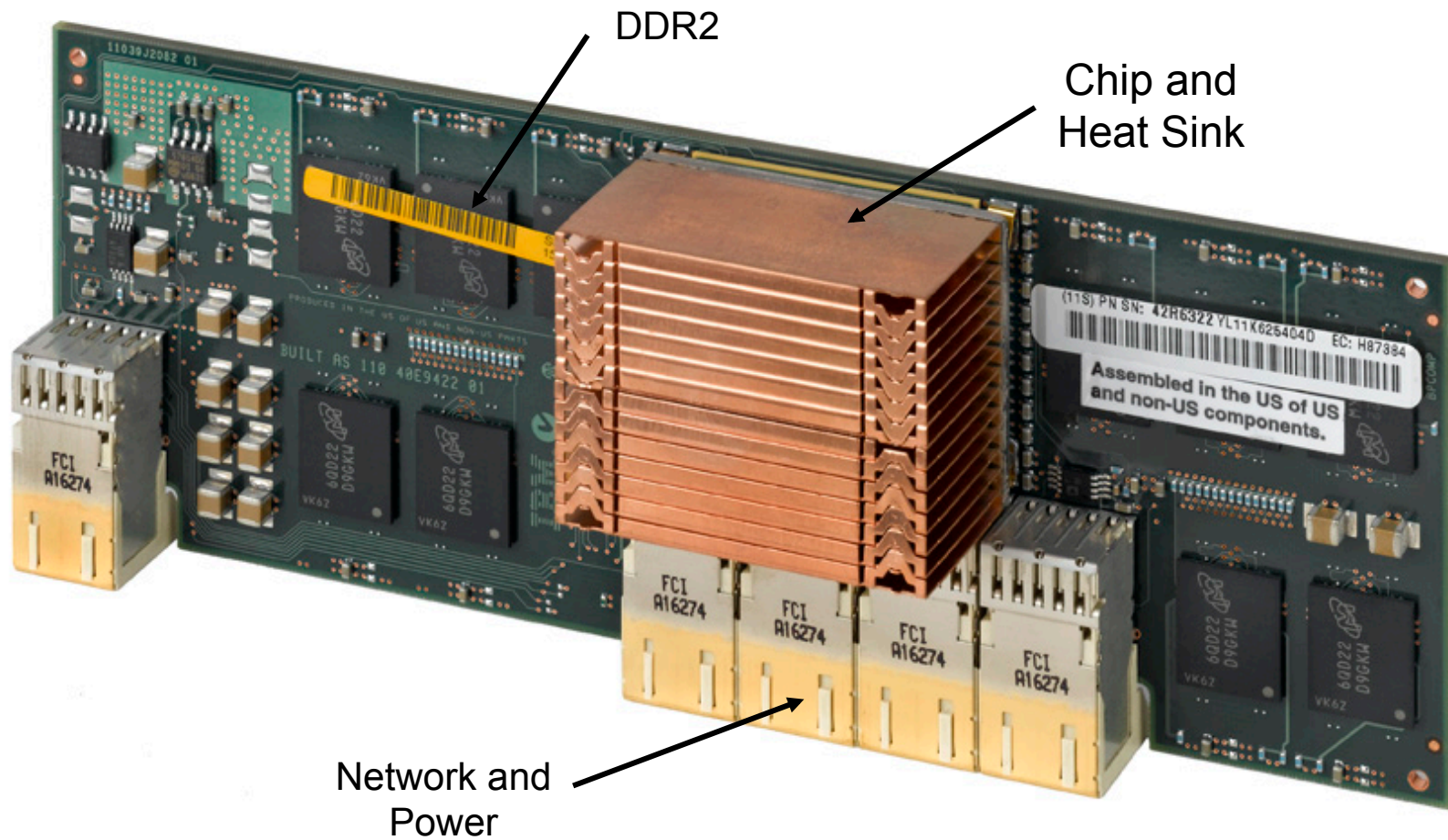
13mmx13mm
90 nm process
208M transistors
88M in eDRAM



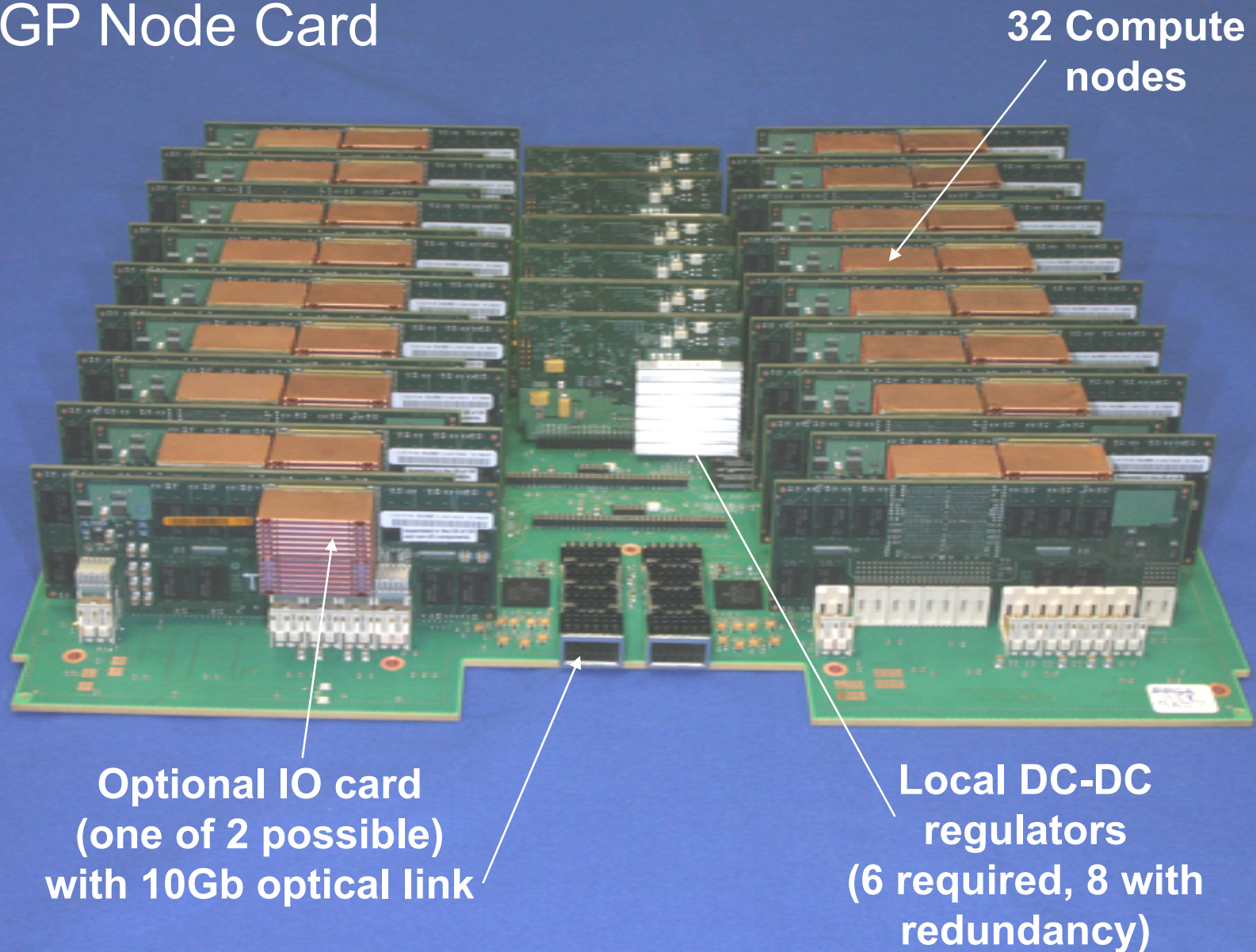
BG/P Compute Card



BG/P Compute Card



BGP Node Card

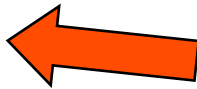


32 Compute Nodes

128 cores

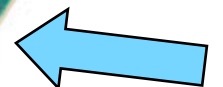
Hottest ASIC Tj
80°C@24W, 55°C@15W

Outlet Air



Max +10°C

Inlet Air



min 2.5m/s
max 17°C

Hottest DRAM
Tcase 75°C@0.3W

Optional IO card
(1 of 2 possible)

10Gb Ethernet

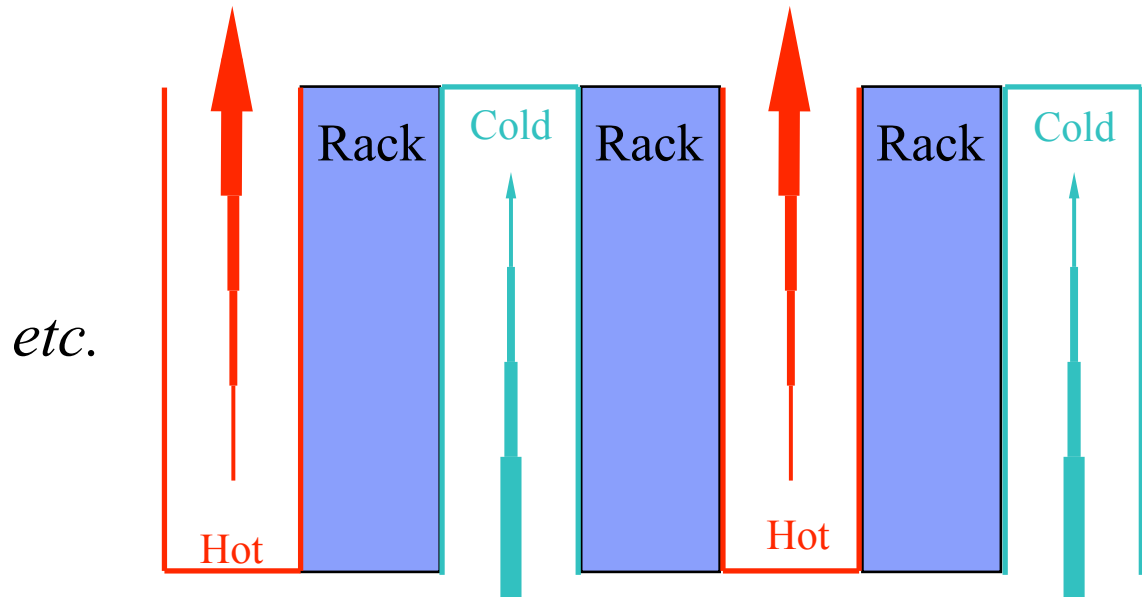
Local 48V input DC-DC regulators
5+1, 3+1 with redundancy. Vicor
technology, tcase 60°C@120A

First BG/P Rack



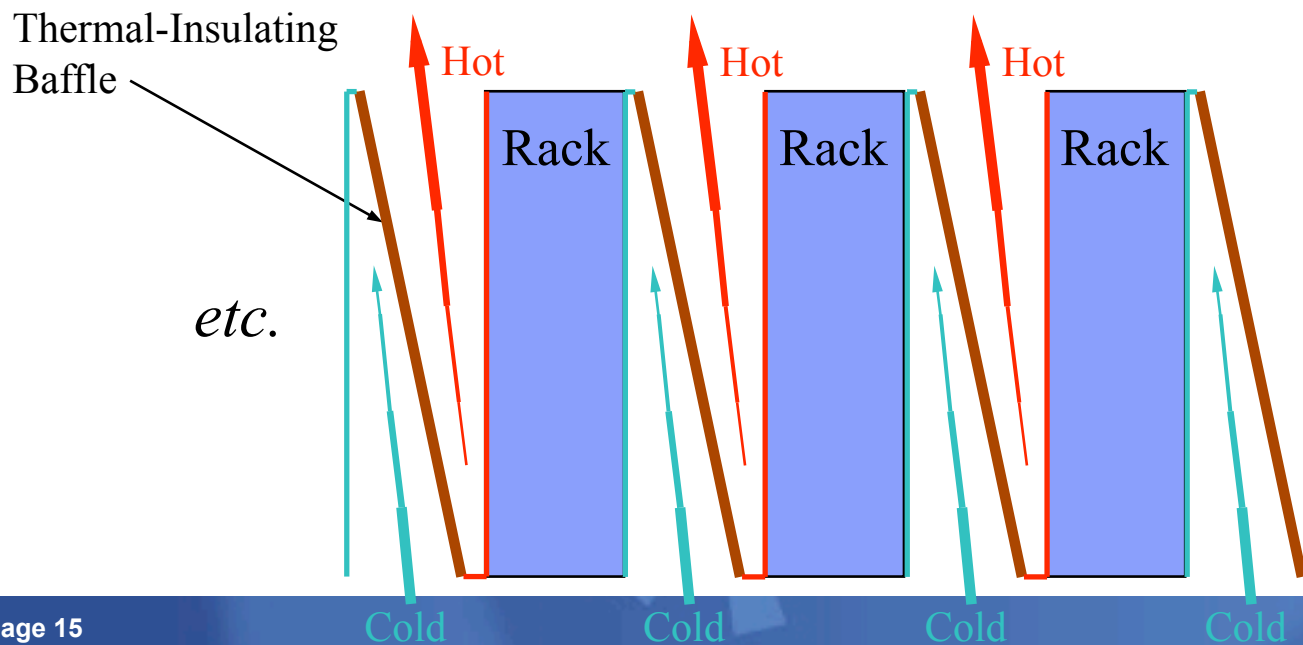
First 8 racks of BG/P: Covers removed





**(a) Prior Art:
Segregated,
Non-Tapered
Plenums**

(Plenum Width Same
Regardless of Flow Rate)



**(b) Invention:
Integrated,
Tapered
Plenums**

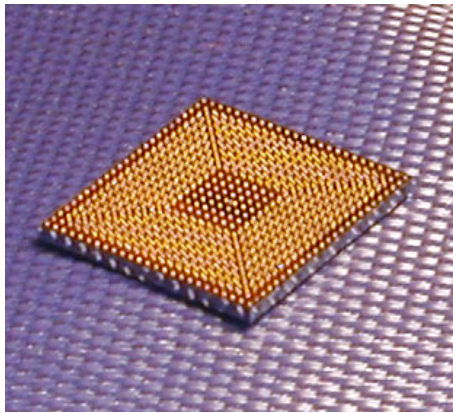
(Plenum Width
Larger where Flow
Rate is Greater)

Shawn Hall 4-3-02
02-04-03 Angled Plenums

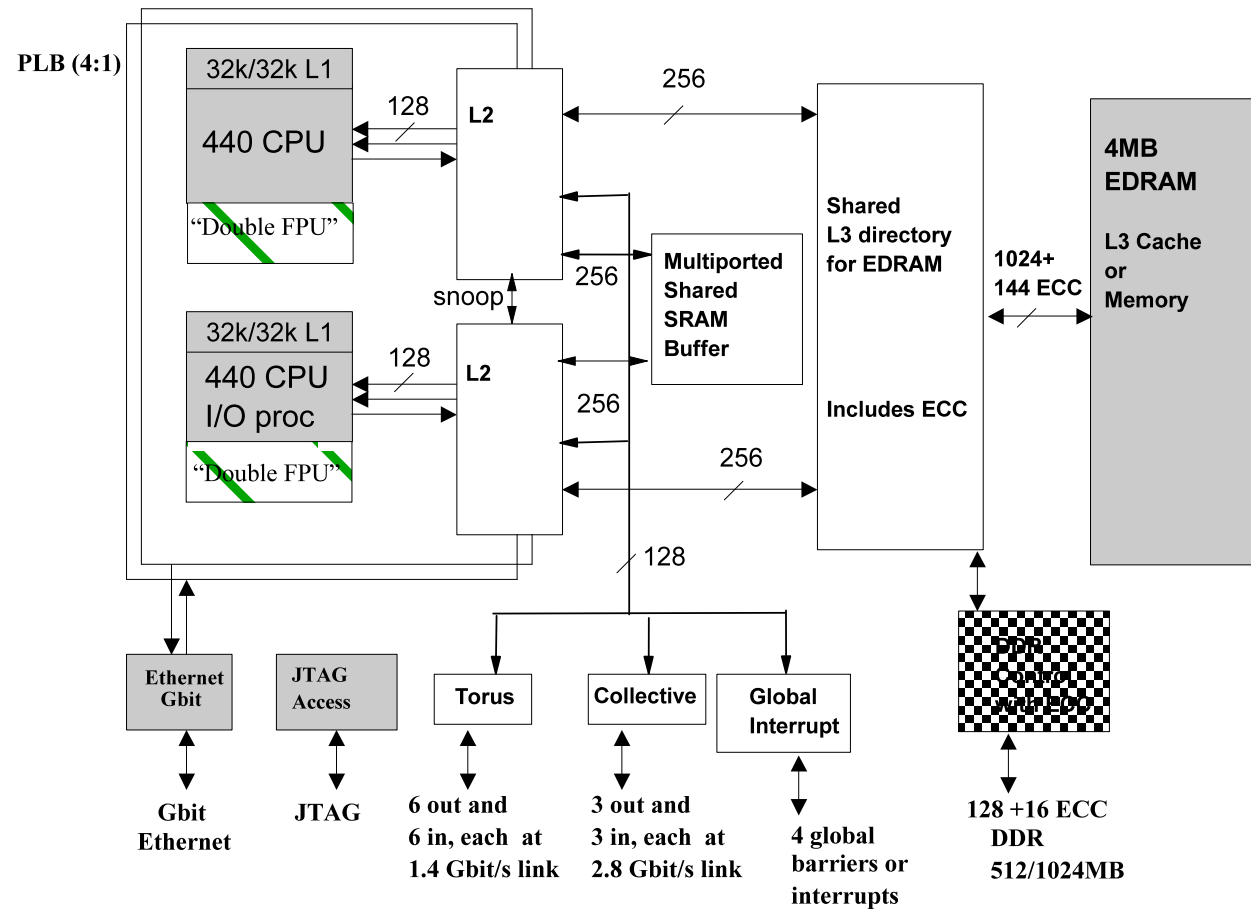
IBM Blue Gene/P



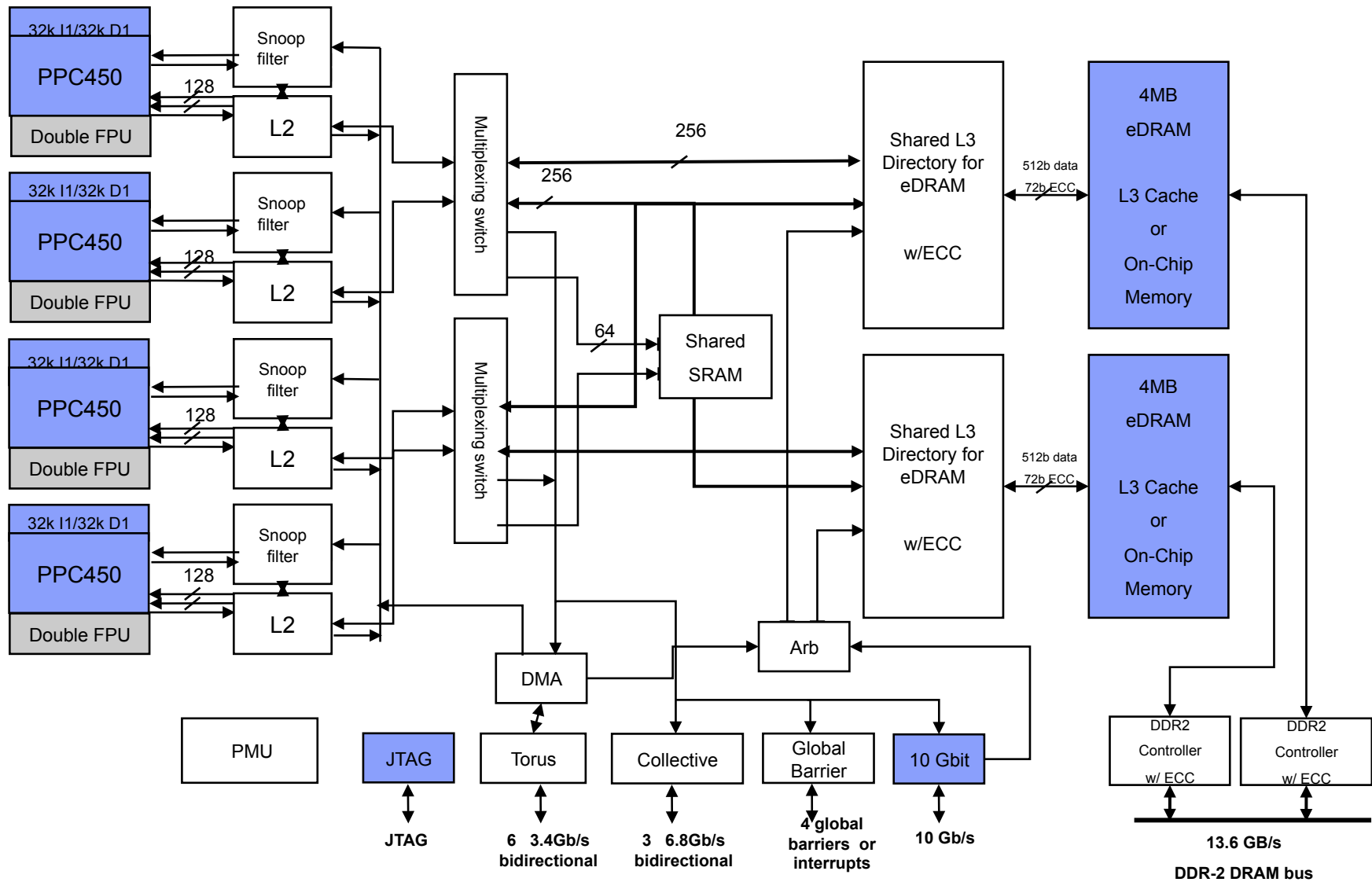
Blue Gene/L ASIC



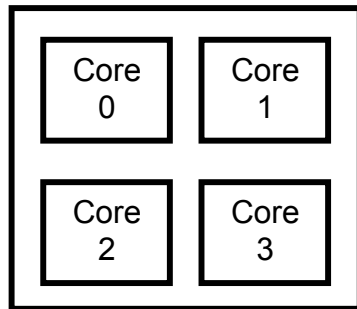
- IBM CU-11, 0.13 μm
- 11 x 11 mm die size
- 25 x 32 mm CBGA
- 474 pins, 328 signal
- 1.5/2.5 Volt



Blue Gene/P ASIC

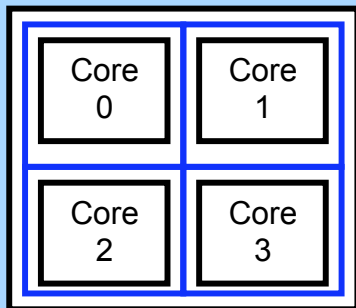


Execution Modes in BG/P

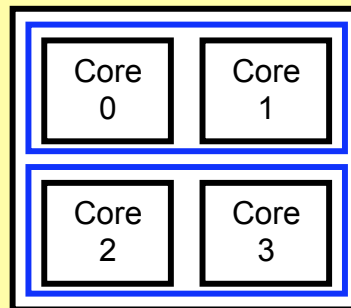


Hardware Elements Black
Software Abstractions Blue

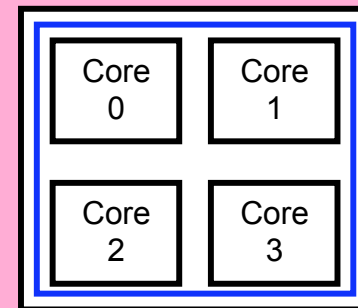
Quad Mode (VNM)
4 Processes
1 Thread/Process



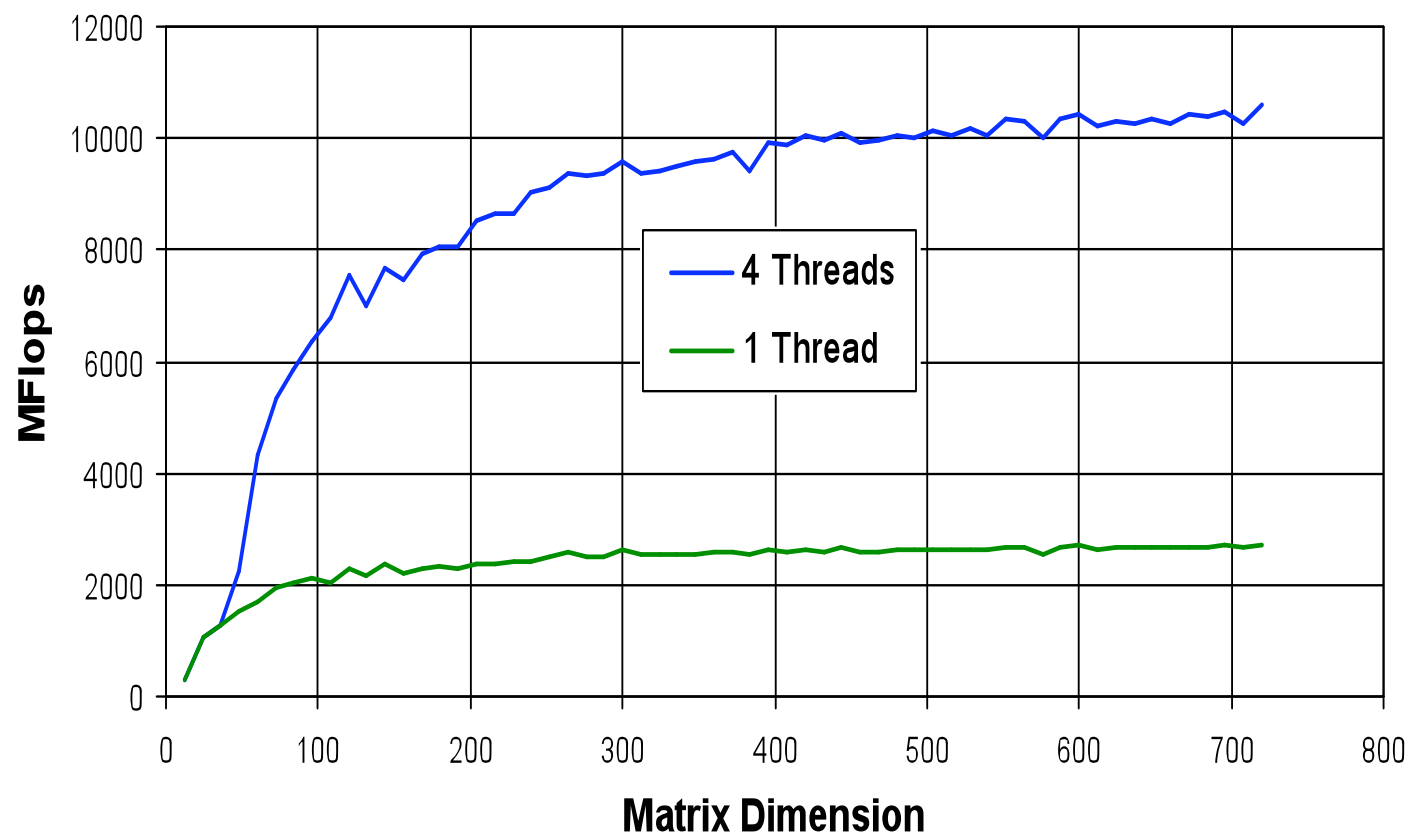
Dual Mode
2 Processes
1-2 Threads/Process



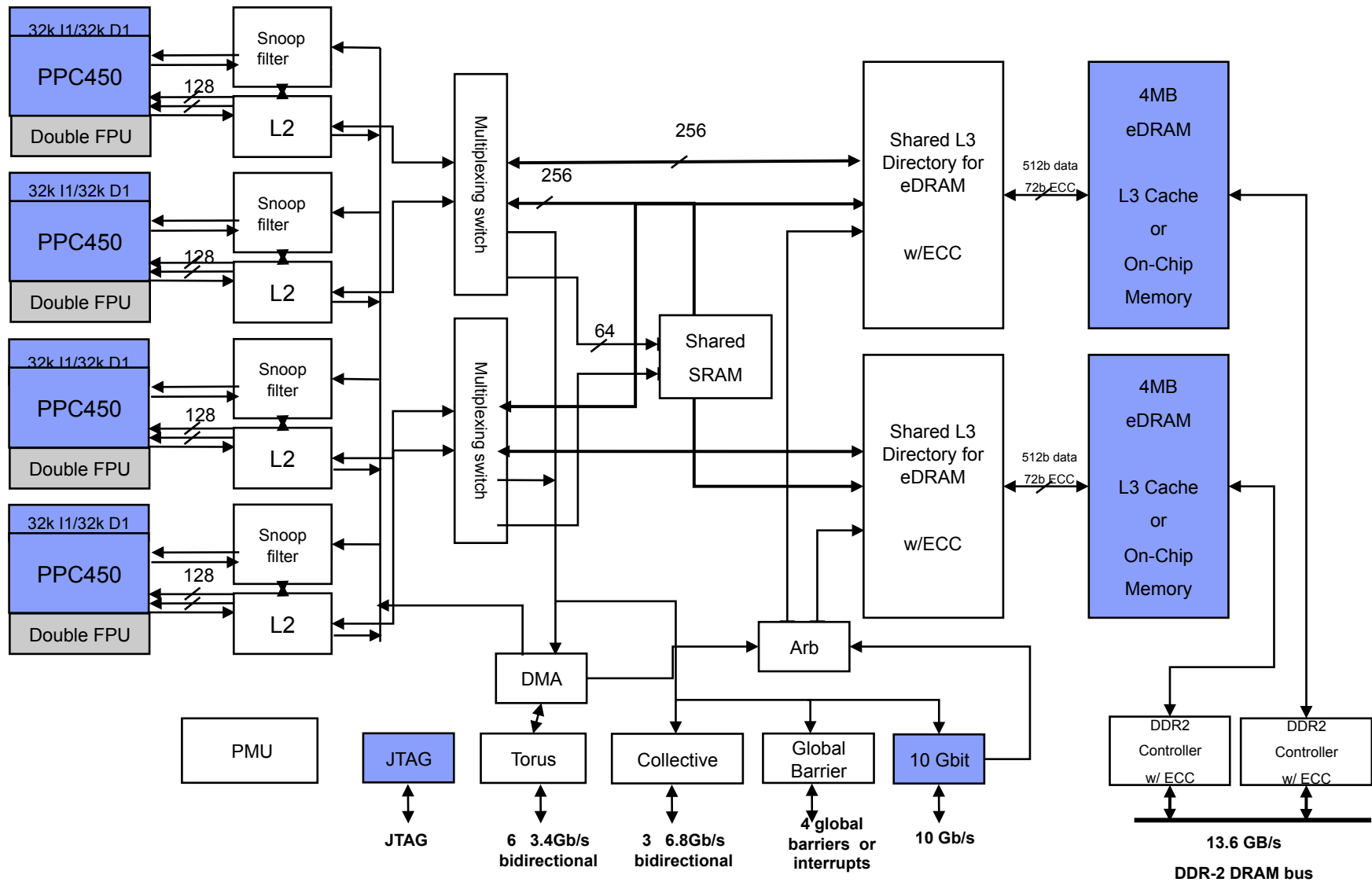
SMP Mode
1 Process
1-4 Threads/Process



BG/P DGEMM (ESSL) One Node



Blue Gene/P ASIC



L1 Cache

— Architecture

- 32KB Instruction-cache, 32KB Data-cache per core
- 32Byte lines, 64 way-set-associative, 16 sets
- Round-robin replacement
- Write-through mode
- No write allocation

— Performance

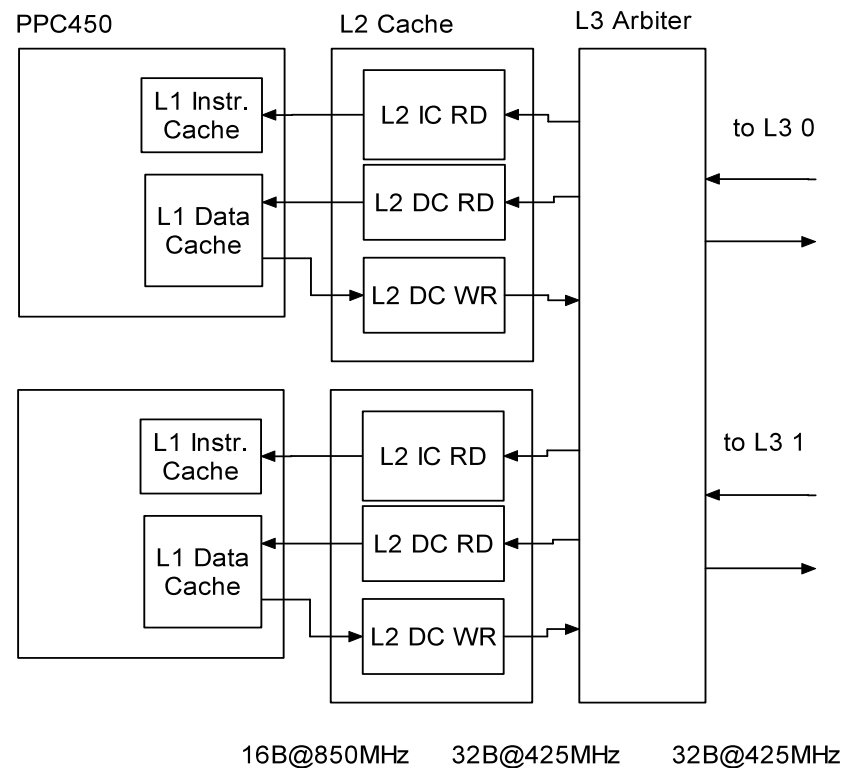
- L1 load hit => 8Bytes/cycle, 4 cycle latency (floating point)
- L1 load miss, L2 hit => 4.6Bytes/cycle, 12 cycle latency
- Store: Write through, limited by external logic to about one request every 2.9 cycles (about 5.5Bytes/cycle peak)

L2 Cache

Three independent ports:
Instruction read
Data Read = prefetch
Data write

L3 Arbiter

Switch function
1 read and 1 write
switched to each L3
every 425MHz cycle



L2 Data Cache Prefetch Unit

Prefetch engine

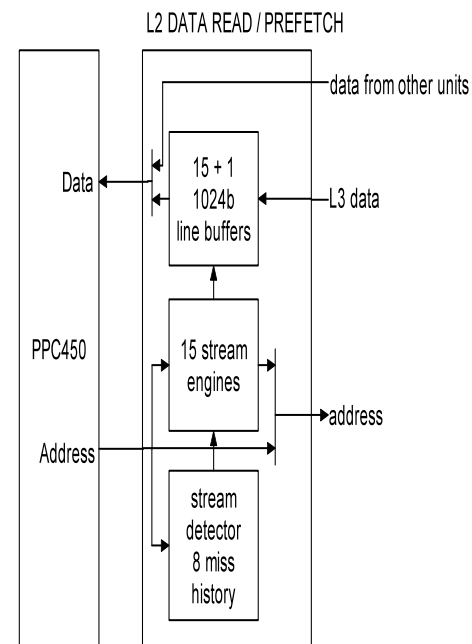
128Byte line prefetch

15 entry fully associative prefetch buffer

1 or 2 line deep prefetching

4.6Bytes/cycle, 12 cycle latency

Supports ~ 7 streams

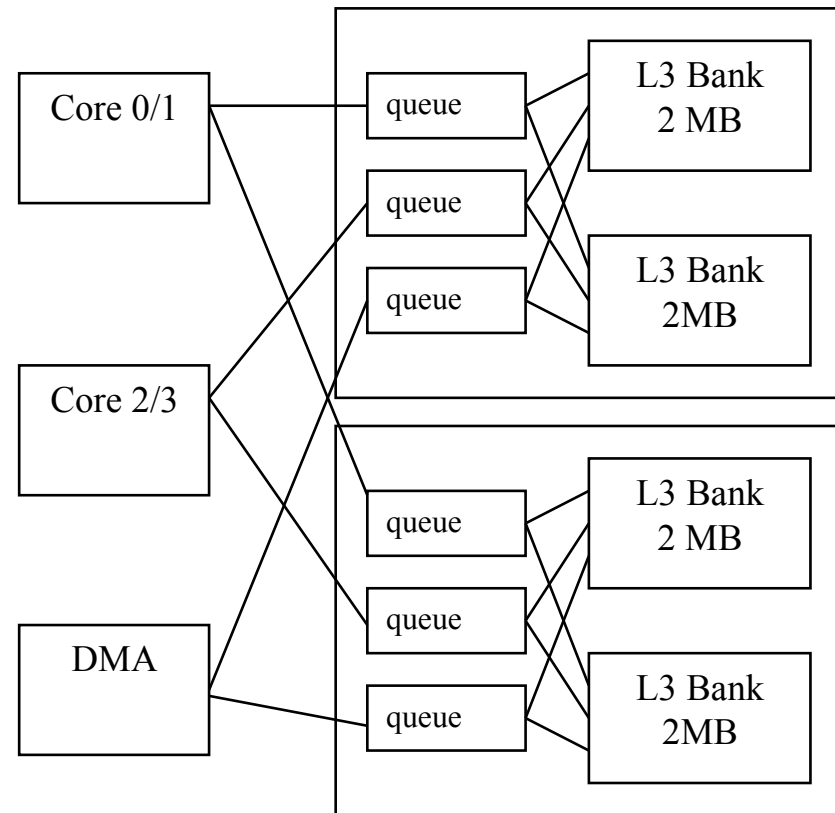


L3 Cache

4 x 2 MB embedded DRAM
banks per node (8MB total),
each containing:

L3 directory

15 entry 128B-wide write
combining buffer



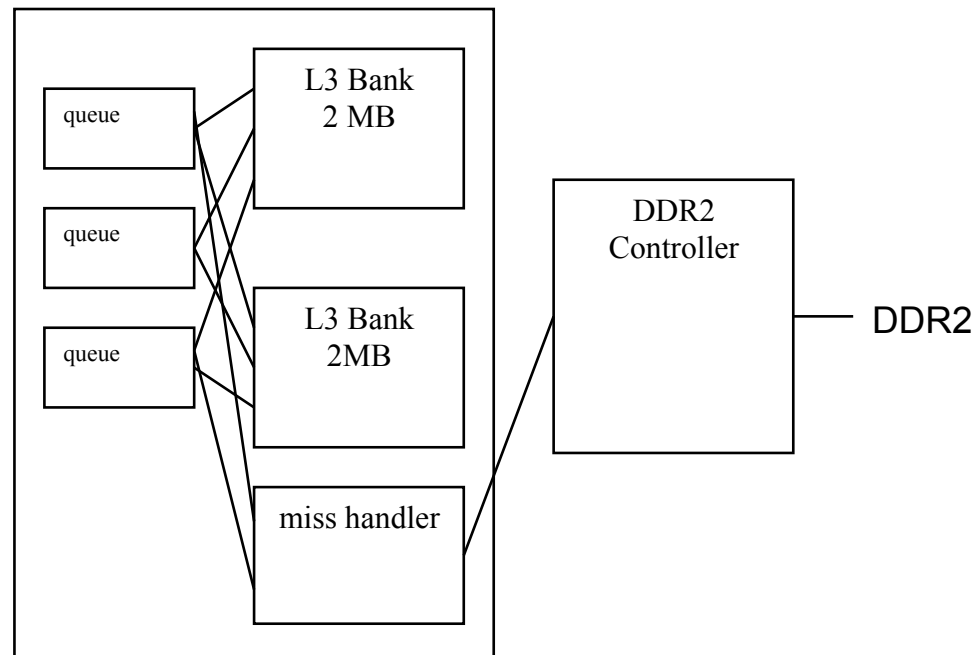
Memory Controllers on Chip

20 8b-wide DDR2
modules per controller

Controllers are on the
chip

4 module-internal banks

Command reordering
based on bank
availability



Memory System Bottlenecks

L2 – L3 switch

- Not a full core to L3 bank crossbar

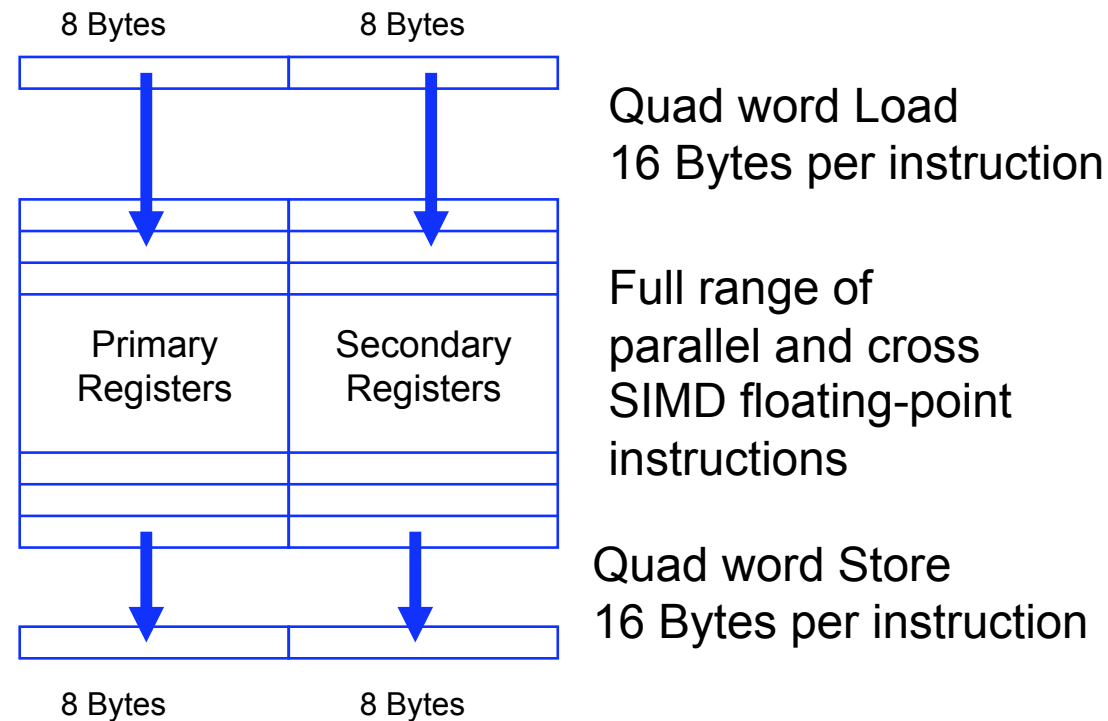
- Request rate and bandwidth are limited if two cores of one dual processor group access the same L3 cache bank

Banking for DDR2

- 4 banks on 512Mb DDR modules

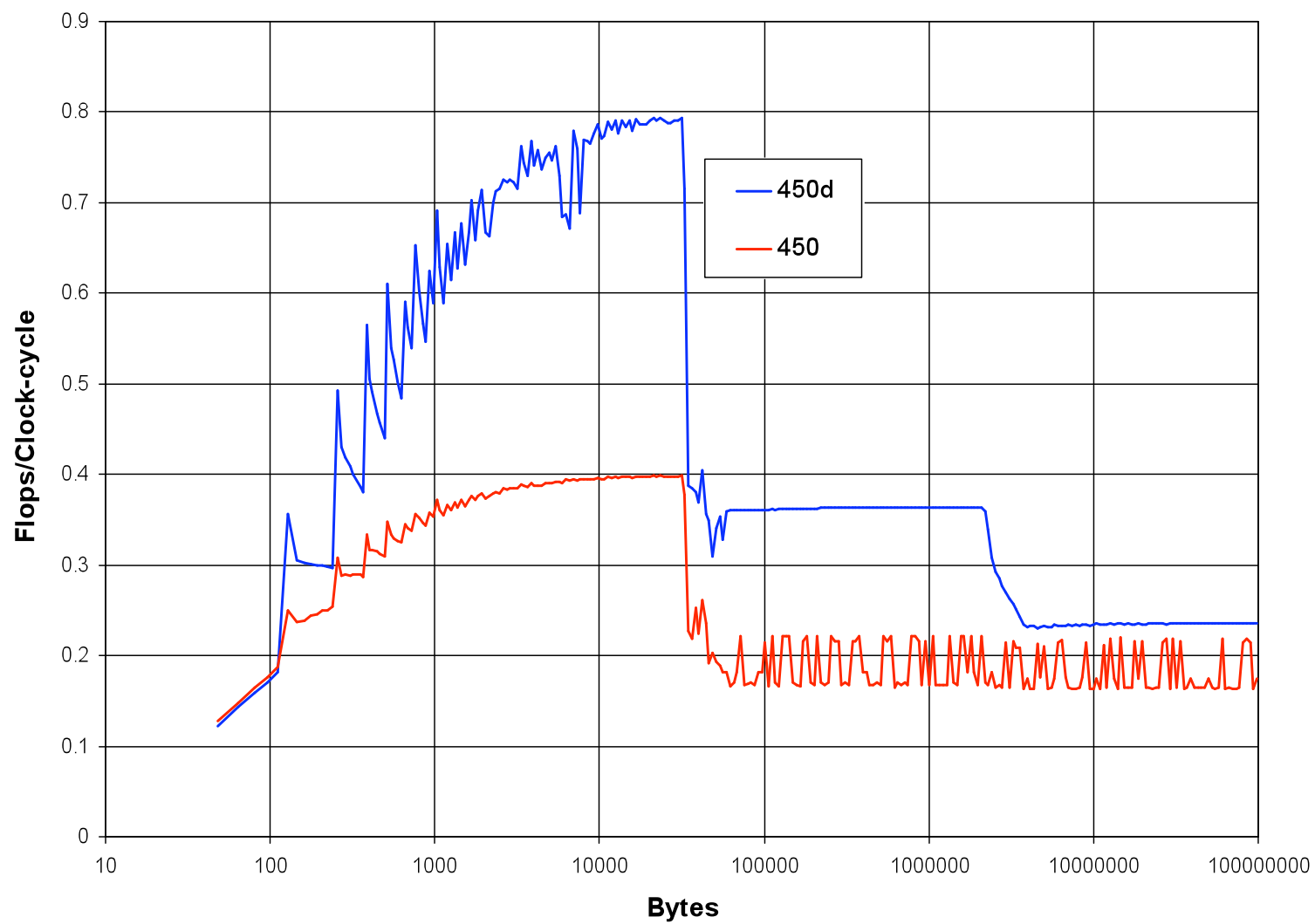
- Peak bandwidth only achievable if accessing 3 other banks before accessing the same bank again

Double Hummer Floating Point Unit



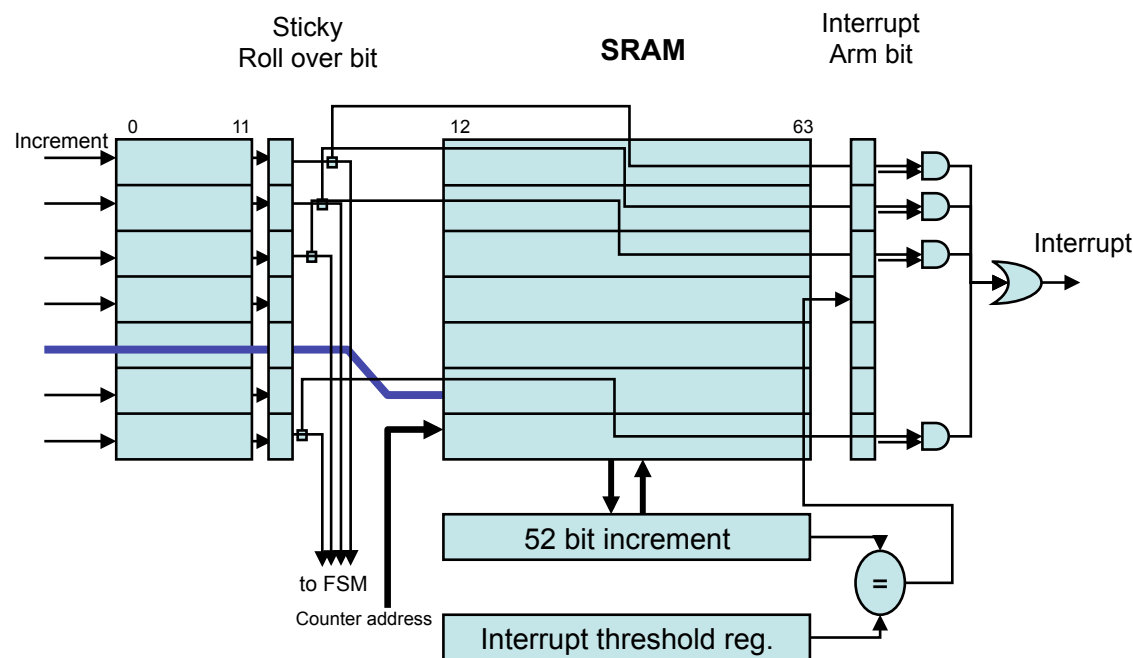
Quad word load/store operations
require data aligned on 16-Byte
boundaries.

BGP Daxpy Performance



Performance Monitor Architecture

- Novel hybrid counter architecture
 - High density and low power using SRAM design
- 256 counters with 64bits resolution
 - Fast interrupt trigger with configurable threshold
 - Performance analysis is key to achieving full system potential



Performance Monitor Features

- Counters for core events
 - loads, stores, floating-point operations (flops)
- Counters for the memory subsystem
 - cache misses, DDR traffic, prefetch info, etc.
- Counters for the network interfaces
 - torus traffic, collective network, DMA, ...
- Counts are tied to hardware elements
 - counts are for cores or nodes, not processes or threads

Blue Gene/P Interconnection Networks



3 Dimensional Torus

- Interconnects all compute nodes
- Virtual cut-through hardware routing
- 3.4 Gb/s on all 12 node links (5.1 GB/s per node)
- 0.5 μ s latency between nearest neighbors, 5 μ s to the farthest
- MPI: 3 μ s latency for one hop, 10 μ s to the farthest
- Communications backbone for point-to-point

Collective Network

- One-to-all broadcast functionality
- Reduction operations for integers and doubles
- 6.8 Gb/s of bandwidth per link per direction
- Latency of one way tree traversal 1.3 μ s, MPI 5 μ s
- Interconnects all compute nodes and I/O nodes

Low Latency Global Barrier and Interrupt

- Latency of one way to reach 72K nodes 0.65 μ s, MPI 1.6 μ s

Blue Gene/P Torus Network

Logic Unchanged from BG/L, *except*

Bandwidth

BG/L:	clocked at $\frac{1}{4}$ processor rate	1Byte per 4 cycles
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BG/P:	clocked at $\frac{1}{2}$ processor rate	1Byte per 2 cycles
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With frequency bump from 700 MHz to 850 MHz

BG/P Links are 2.4x faster than BG/L

425 MB/s vs **175** MB/s

Same Network Bandwidth per Flops as BG/L

Minor changes in error reporting

Primary interface is via DMA, rather than cores

Run application in DMA mode, or core mode (not mixed)

Software product stack uses DMA mode

DMA Overview

Rich function DMA offloads data movement to/from torus

Message Types

Direct Put:	copy from source address to destination address
Remote Get:	remote copy from target node address to source
Memory Fifo:	place packets in Fifo (in memory) on destination

Arbitrary offsets

Can do memcpy within a node

Basic Constructs:

- Injection and Reception Fifos

- Injection and Reception “Counters”

- Organized into 4 “Groups” (could have one group/core)

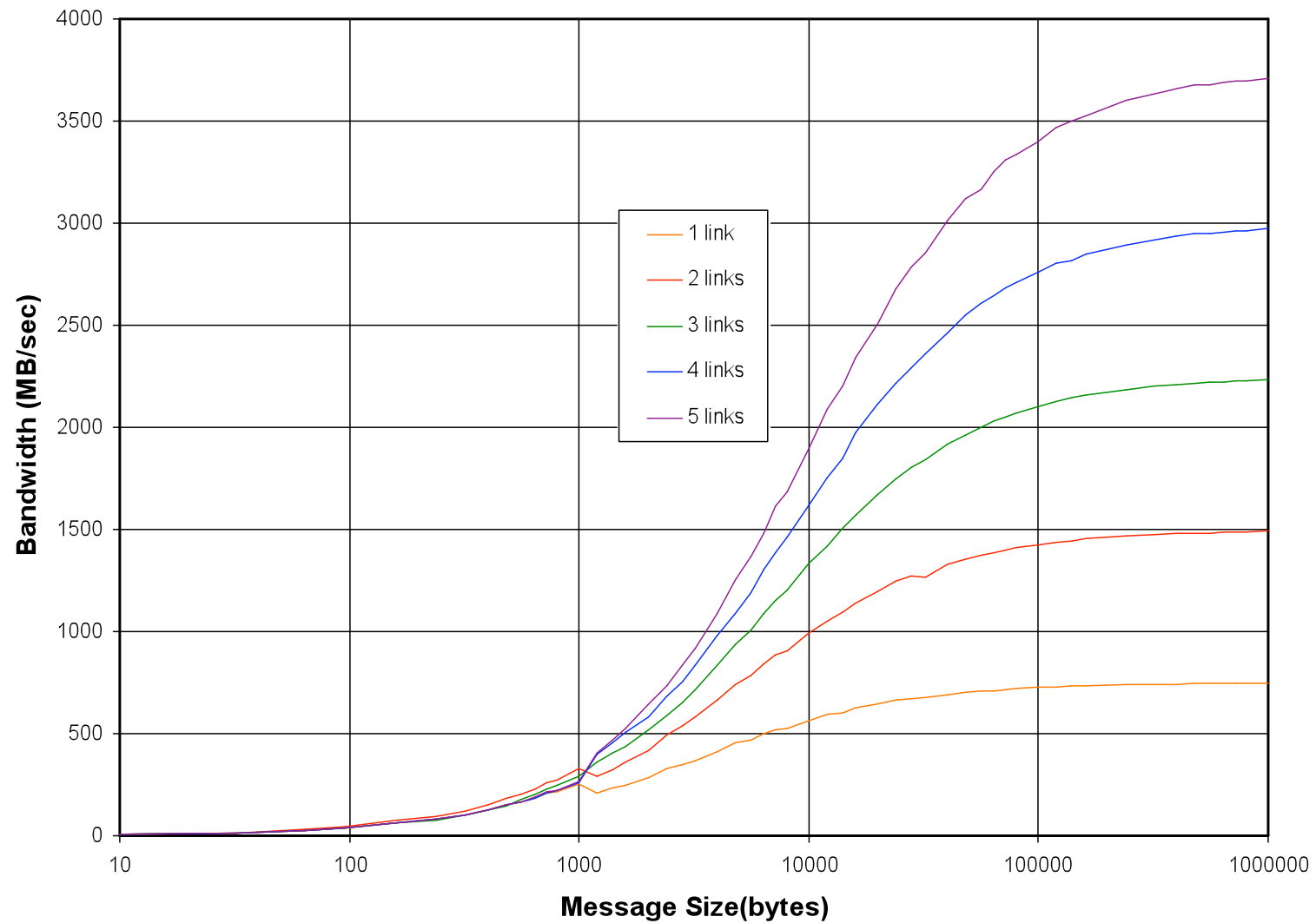
- Message Descriptors (32 B)

 - Includes destination, message type and length, payload pointer

DMA uses physical addresses

- Efficient SPI programming interface for Virtual to Real translation

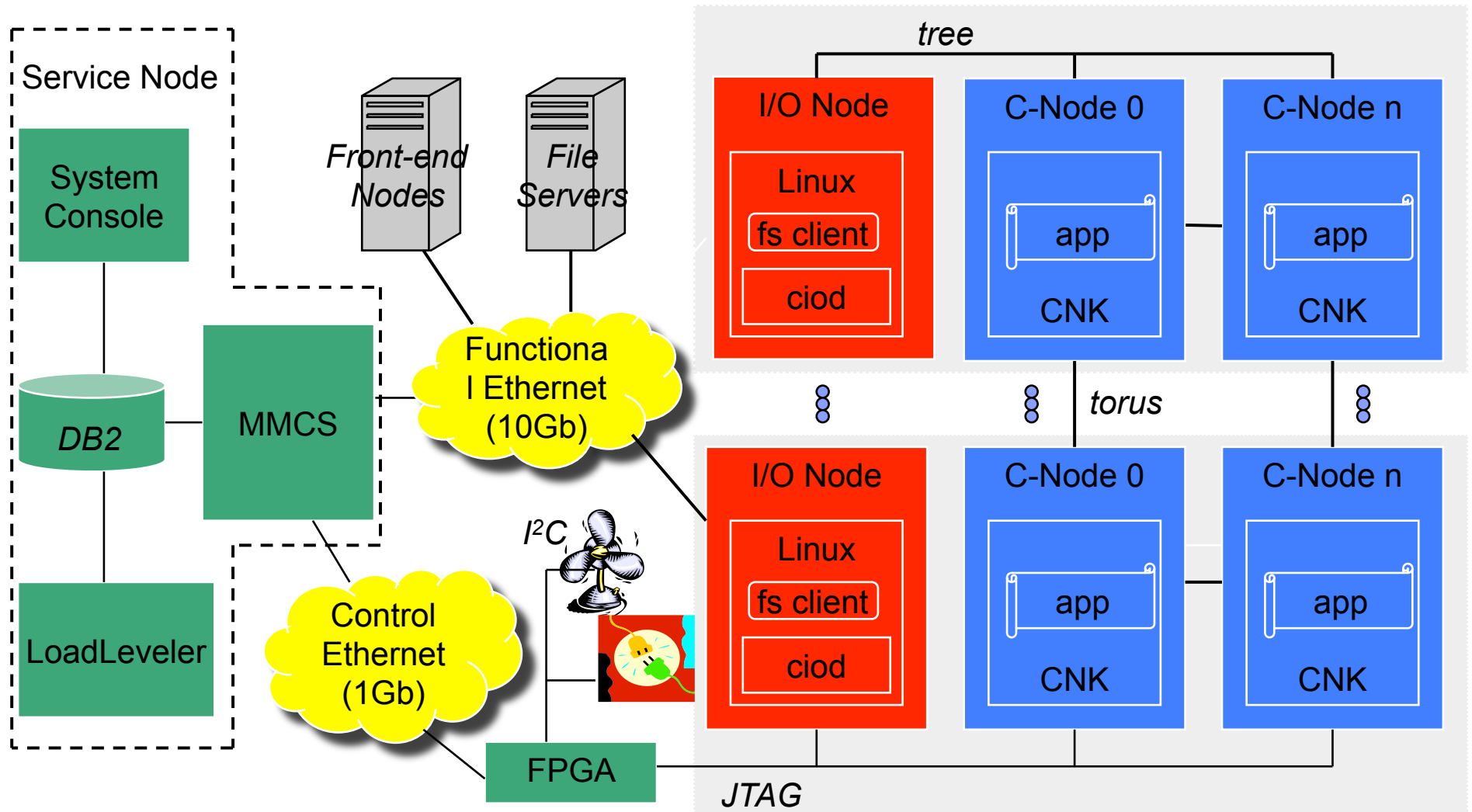
BGP Exchange Bandwidth



Blue Gene/P key architectural improvements over BG/L

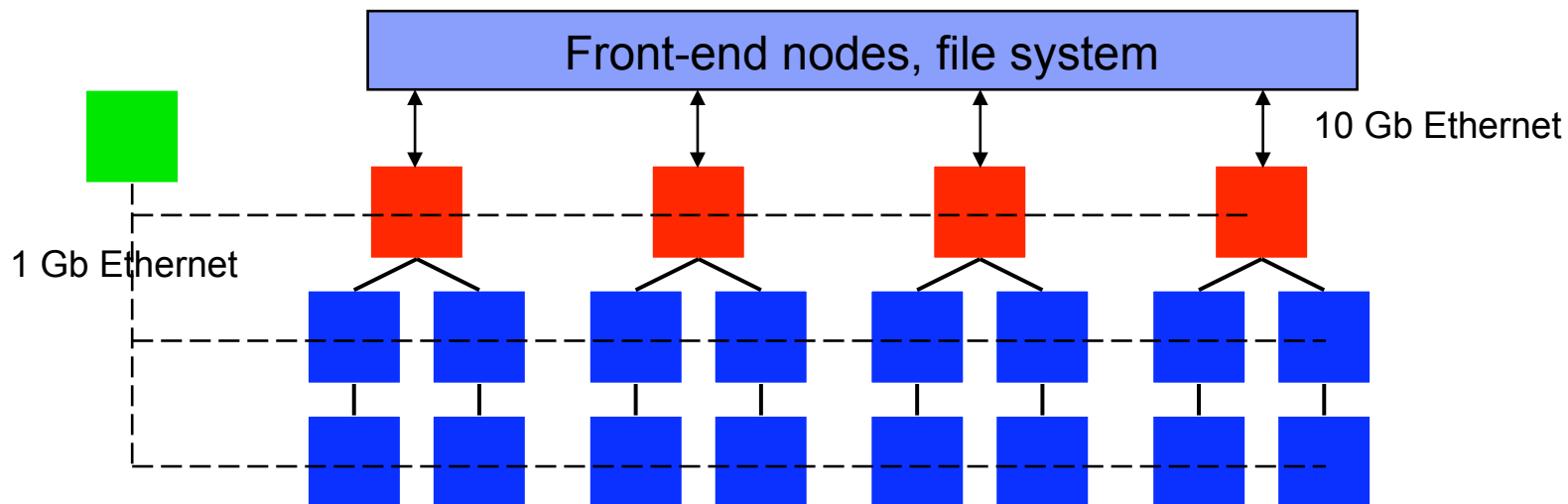
Property		Blue Gene/L	Blue Gene/P
Node Properties	Processor cores/chip	two 440 PowerPC	four 450 PowerPC
	Processor Frequency	0.7GHz	0.85GHz
	Coherency	Software managed	SMP with snoop filtering
	L1 Cache (private)	32KB I + 32KB D/proc.	32KB I + 32KB D/proc.
	L2 Cache (private)	15 line buffers	15 line buffers
	L3 Cache size (shared)	4MB	8MB
	Main Store	512 MB and 1 GB DDR	2 GB DDR2
	Main Store Bandwidth	5.6 GB/s (16B wide)	13.6 GB/s (2*16B wide)
	Peak Performance	5.6 GFLOPs/node	13.6 GFLOPs/node
Torus Network	Aggregate Bandwidth	6*2*175 MB/s=2.1 GB/s	6*2*425 MB/s= 5.1 GB/s
	Hardware Latency (Nearest Neighbor)	<1 μ s	<1 μ s
Collective Network	Aggregate Bandwidth	3*2*350 MB/s=2.1 GB/s	3*2*0.85 GB/s=5.1 GB/s
Performance Monitors	Counters	48	256
	Counter resolution (bits)	32b	64b
GFLOPS/Watt		0.22	0.33

Blue Gene/P System Architecture



Blue Gene Software Hierarchical Organization

- **Compute nodes** dedicated to running user application, and almost nothing else - simple compute node kernel (CNK)
- **I/O nodes** run Linux and provide a more complete range of OS services – files, sockets, process launch, signaling, debugging, and termination
- **Service node** performs system management services (e.g., partitioning, heart beating, monitoring errors) - transparent to application software

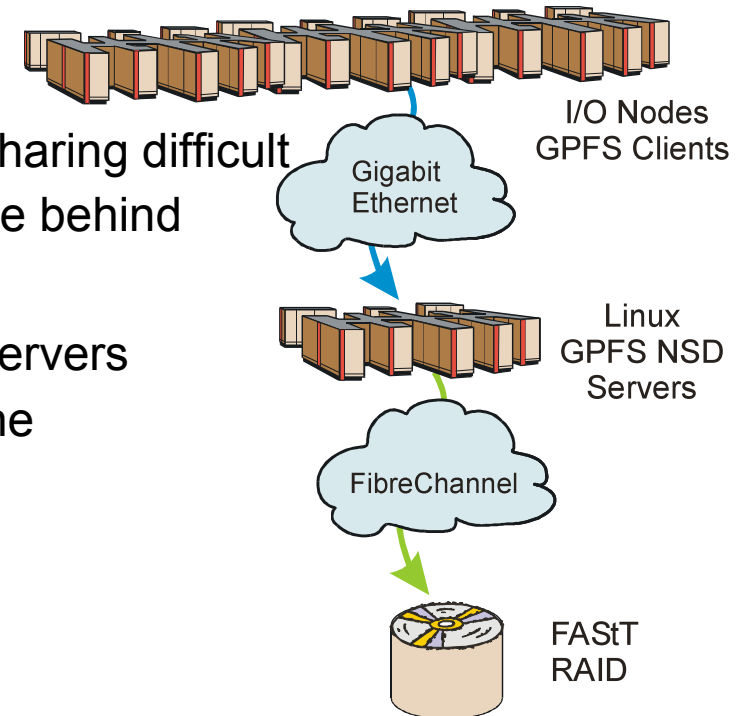


Programming models and development environment

- Familiar methods
 - SPMD model - Fortran, C, C++ with MPI (MPI1 + subset of MPI2)
 - Full language support with IBM XL and GNU compilers
 - Automatic SIMD FPU exploitation (limited)
 - Linux development environment
 - User interacts with system through front-end nodes running Linux – compilation, job submission, debugging
 - Compute Node Kernel provides look and feel of a Linux environment
 - POSIX routines (with some restrictions: no fork() or system())
 - BG/P adds pthread support, additional socket support
 - Tools – support for debuggers, MPI tracer, profiler, hardware performance monitors, visualizer (HPC Toolkit), PAPI
- Restrictions (which lead to significant benefits)
 - Space sharing - one parallel job per partition of machine, one thread per core in each compute node
 - Virtual memory is constrained to physical memory size

General Parallel File System (GPFS) for Blue Gene

- Blue Gene can generate enormous I/O demand (disk limited)
 - BG/P IO-rich has 64 10Gb/rack – 80GB/sec
- Serving this kind of demand requires a parallel file system
- NFS for file I/O
 - Limited scalability
 - NFS has no cache consistency, making write sharing difficult
 - Poor performance, not enough read ahead/write behind
- GPFS runs on Blue Gene
 - GPFS clients in Blue Gene call external NSD servers
 - Brings traditional benefits of GPFS to Blue Gene
 - I/O parallelism
 - Cache consistent shared access
 - Aggressive read-ahead, write-behind



BG/P Single Node Software Enhancements

- Compute Node Kernel
 - SMP support – multithreading
 - Enhancements to support broader class of applications on-demand linking, I/O services
 - Support for enhanced virtual node mode with sharing of read-only data
- I/O node kernel
 - Full SMP support
 - Drive 10 Gb ethernet
- Compilers
 - OpenMP support to version 2.5 specification
- Math libraries
 - ESSL, MASS, MASSV – complete functionality

BG/P Multi-node Software Enhancements

- MPI
 - Full MPI2 support, except for dynamic process management
 - MPI-IO optimizations, one-sided communication
 - Exploit DMA for the torus network
 - Exploit multiple processors for single MPI call
 - Higher levels of scalability
 - Adaptive protocol selection
 - Adaptive connection and buffer management
- Control system
 - Enhanced ease of management
 - Higher levels of scalability
 - Proactive handling of failures
- Job scheduler/launch
 - Integration in a grid environment – making Blue Gene available as a managed web service

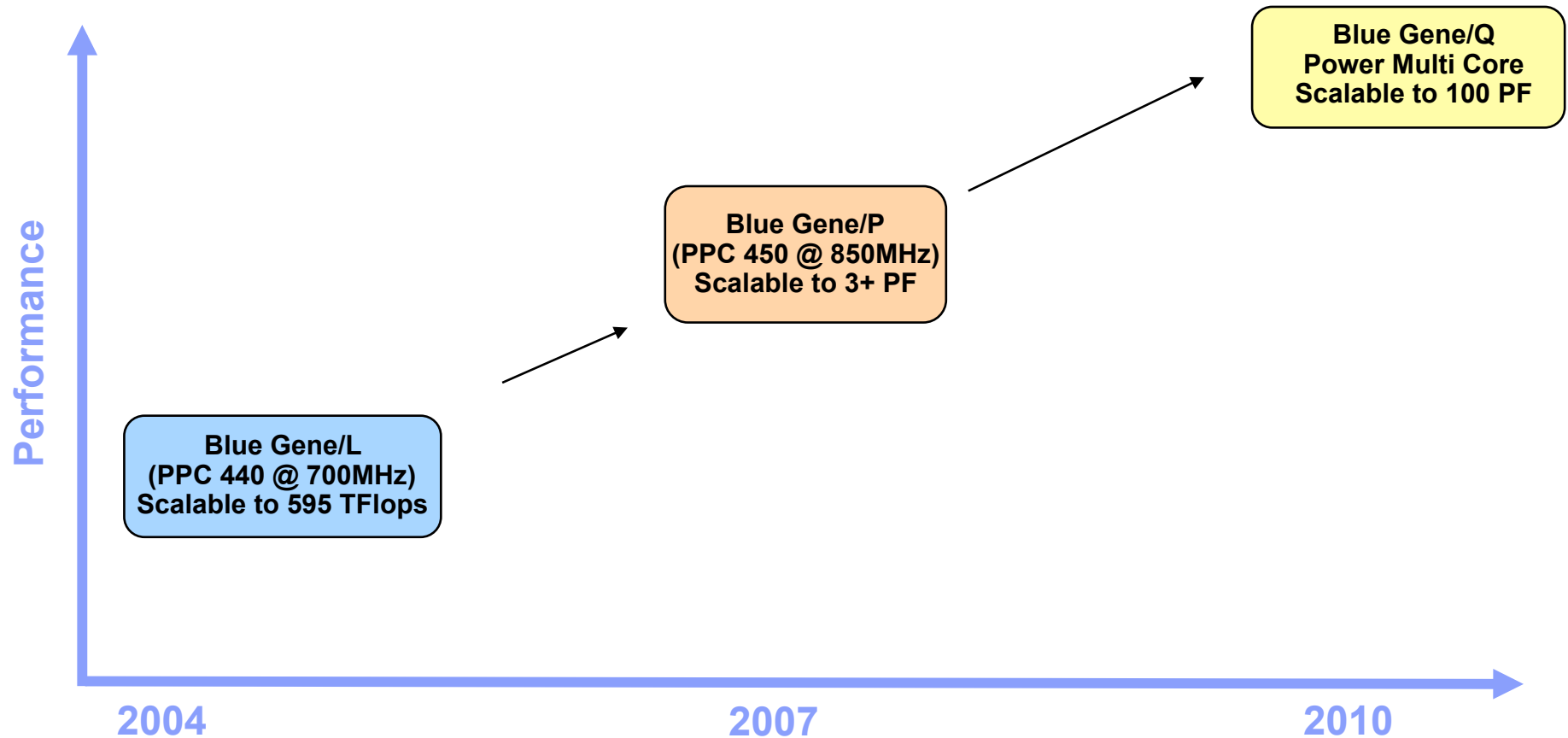
BG/P Linpack Performance

System Size	Nodes	Processors	Performance (TF)	Fraction of Peak
1 Rack	1024	4096	11.10	80%
2 Racks	2048	8192	21.91	79%
4 Racks	4096	16384	43.96	79%
8 Racks	8192	32768	85.98	77%
16 Racks	16384	65536	(Top 500) 167.20 (ANL Dec 13) 171.80	75% 77%

Summary

- Blue Gene/P: Facilitating Extreme Scalability
 - Ultrascale capability computing
 - Provides customer with enough computing resources to help solve grand challenge problems
 - Provide competitive advantages for customers' applications
 - Energy conscious solution supporting green initiatives
 - Familiar open/standards operating environment
 - Simple porting of parallel codes
- Key Solution Highlights
 - Leadership performance, space saving design, low power requirements, high reliability, and easy manageability

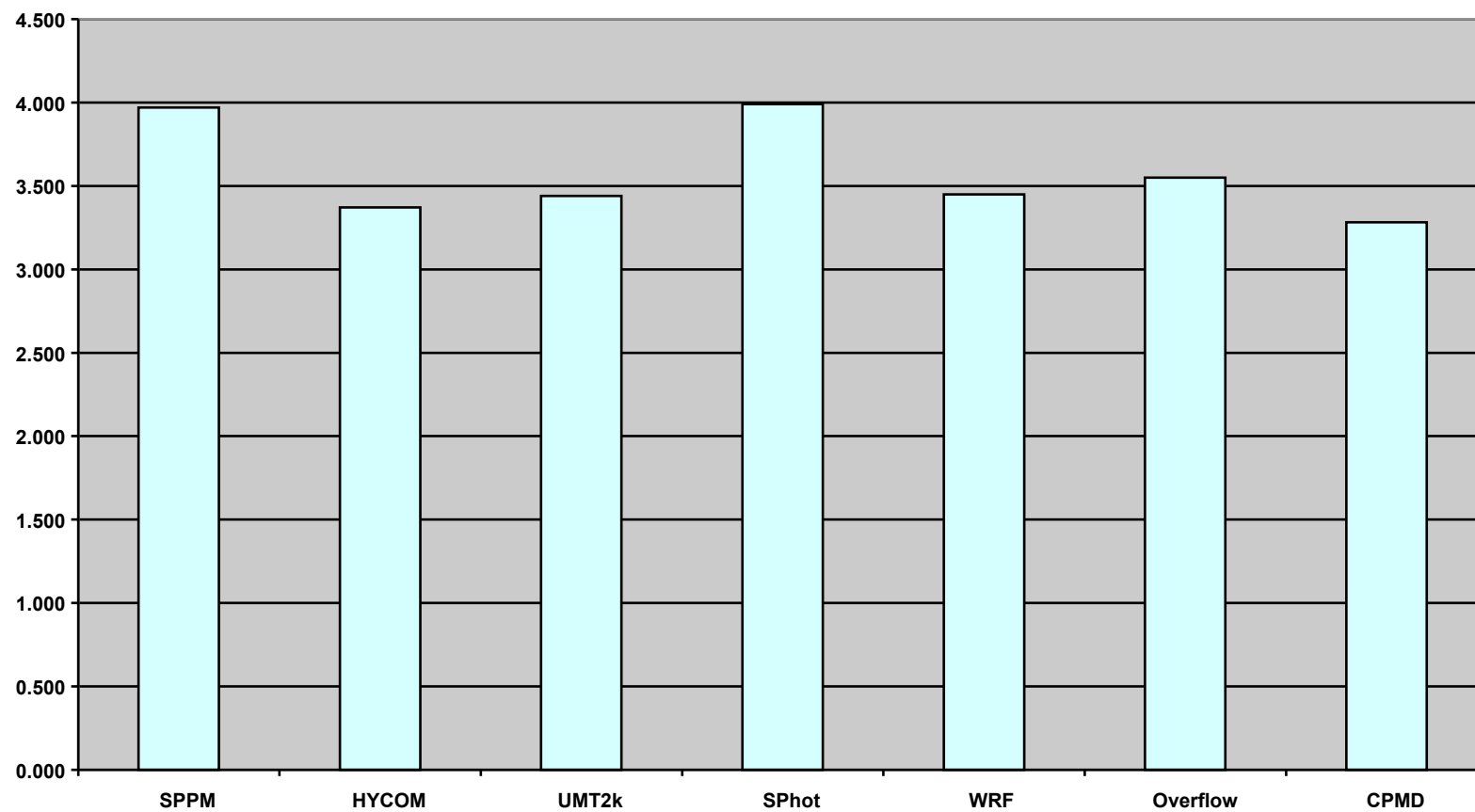
Blue Gene Technology Roadmap



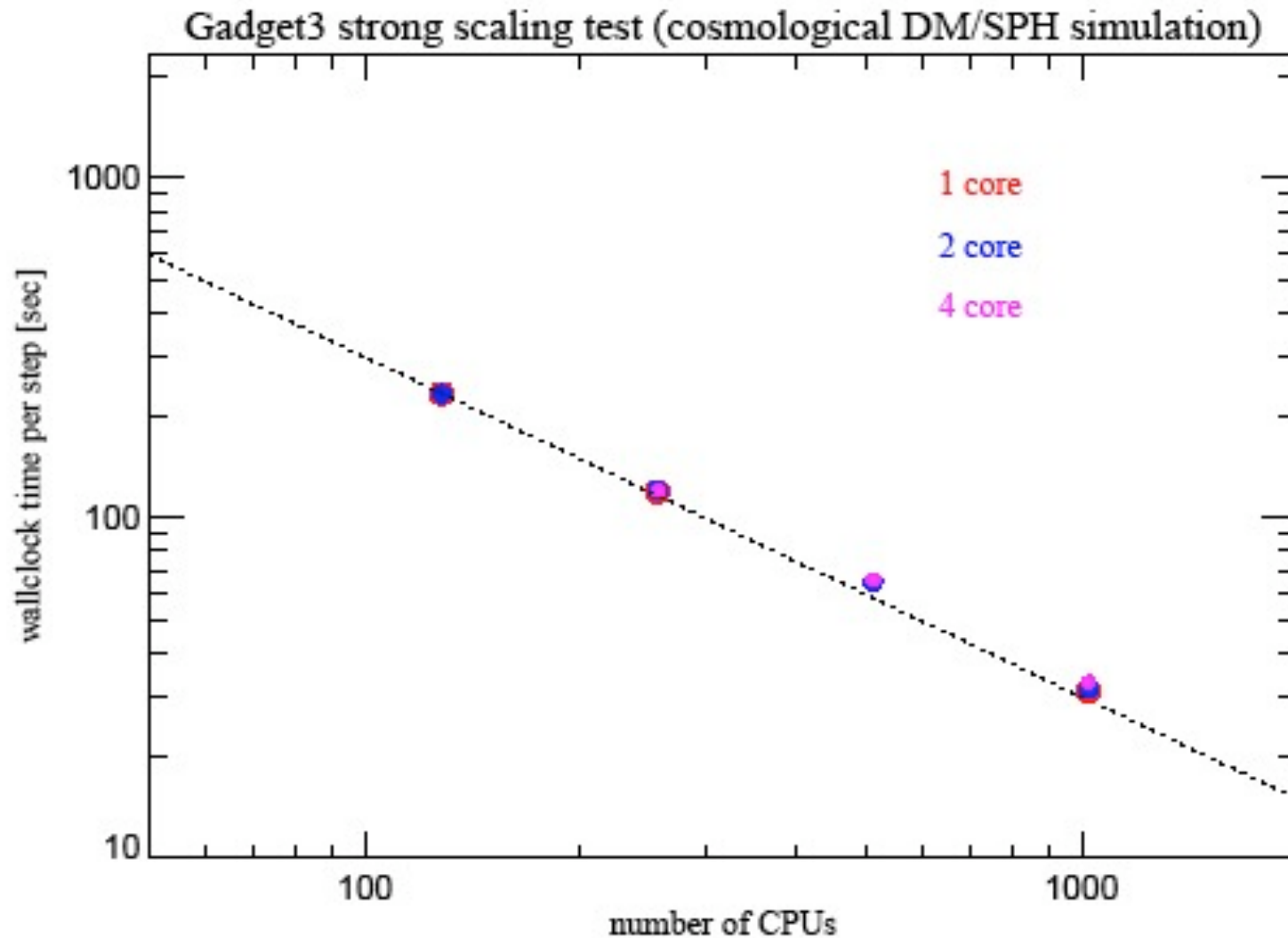
Note: All statements regarding IBM's future direction and intent are subject to change or withdrawal without notice, and represent goals and objectives only.

OpenMP Scaling

Misc codes: OMP Speedup on 4 threads wrt a single thread



Strong Scaling of Applications on BG/P



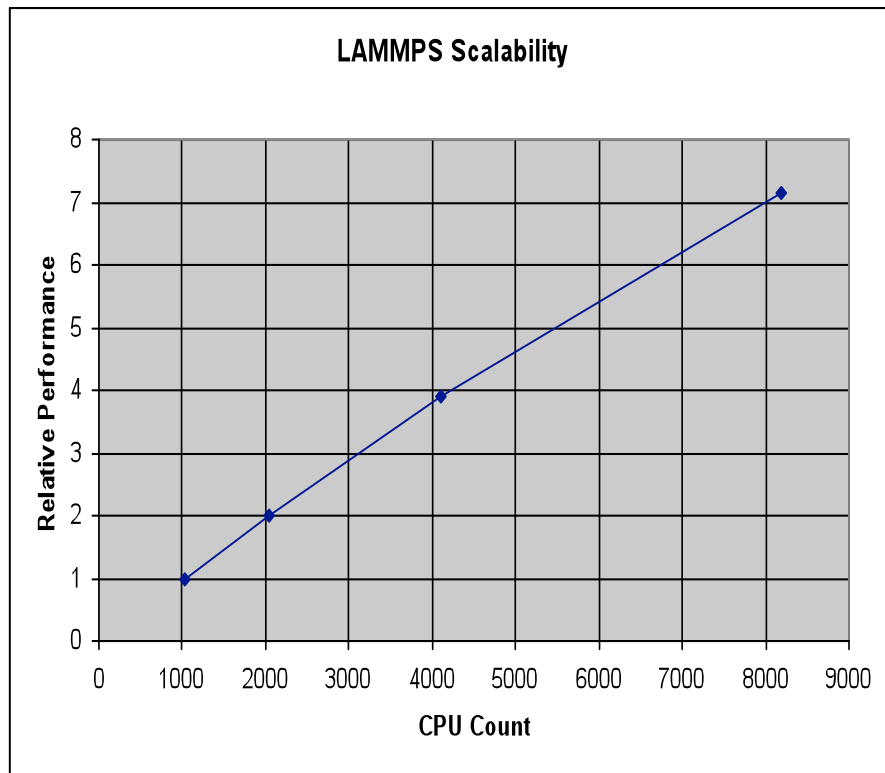
Courtesy
of
Volker
Springel

Max
Planck
Institute
for
Astro-
physics
Garching

Strong Scaling of Applications on BG/P

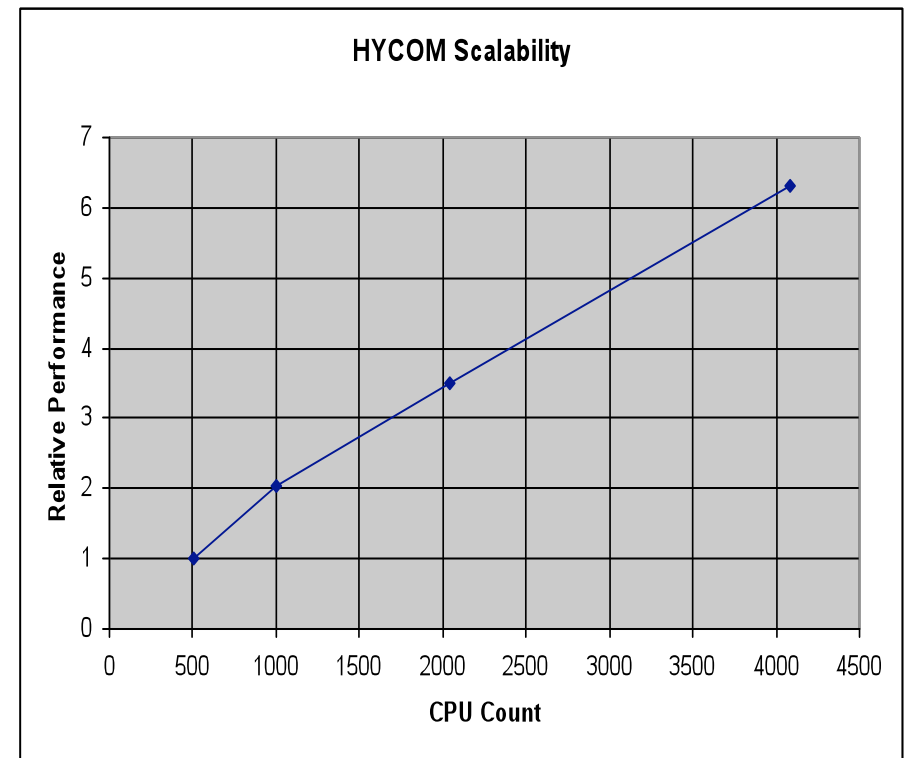
■ LAMMPS

- Molecular Dynamics
- Strong Scaling Problem (fixed problem size)



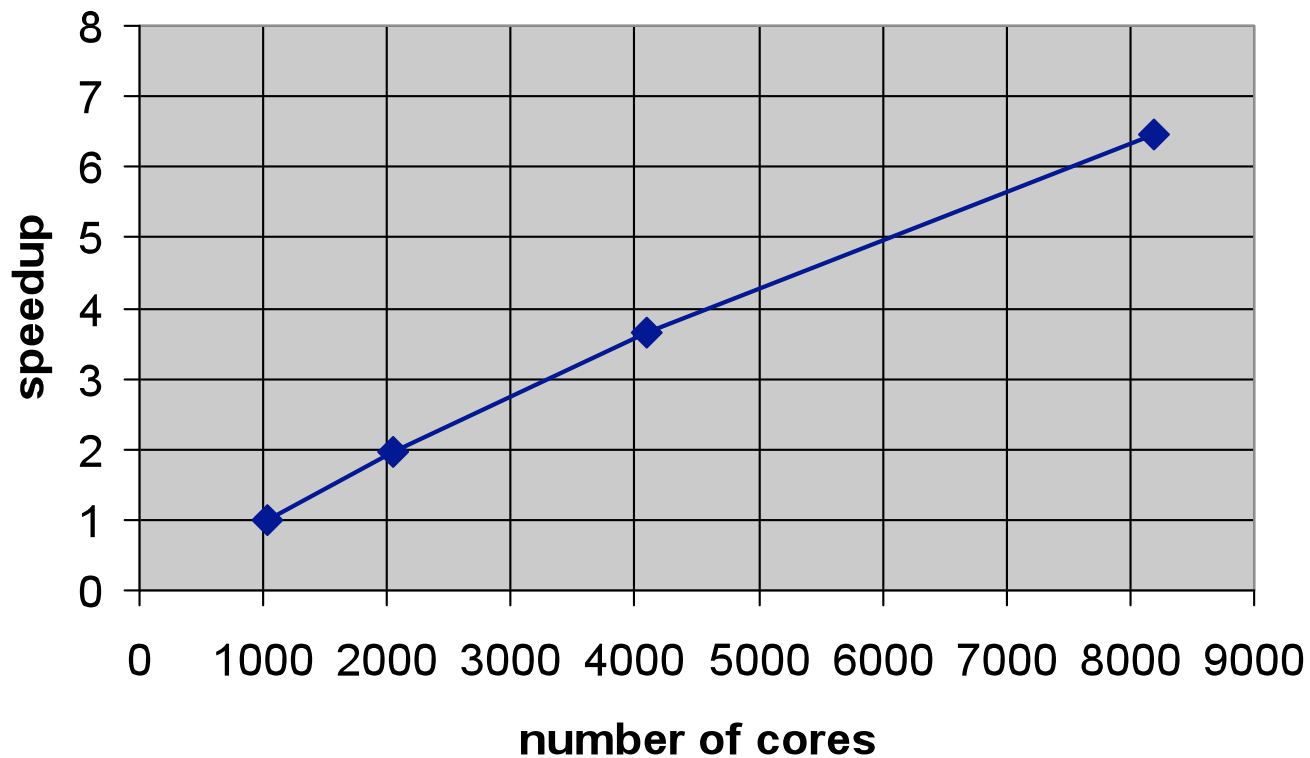
■ HYCOM

- Ocean Model
- Strong Scaling Problem (fixed problem size)



GENE Scaling on BlueGene/P

strong scaling, problem size 0.5 TB
using 4 cores per node



Gyrokinetic
Electromagnetic
Numerical
Experiment
Max Planck